

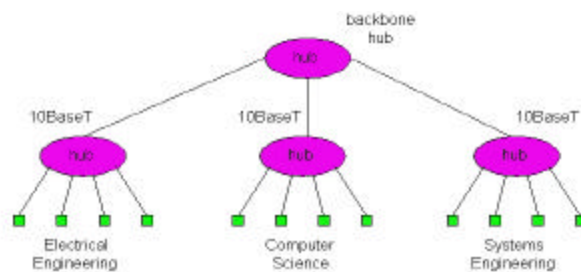
Hubs, Bridges, and Switches (oh my)

- r Used for extending LANs in terms of geographical coverage, number of nodes, administration capabilities, etc.
- r Differ in regards to:
 - m collision domain isolation
 - m layer at which they operate
- r Different than routers
 - m plug and play
 - m don't provide optimal routing of IP packets

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Hubs

- r Physical Layer devices: essentially repeaters operating at bit levels: repeat received bits on one interface to all other interfaces
- r Hubs can be arranged in a hierarchy (or **multi-tier design**), with a **backbone hub** at its top



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Hubs (more)

- r Each connected LAN is referred to as a LAN **segment**
- r Hubs **do not isolate** collision domains: a node may collide with any node residing at any segment in the LAN

- r Hub Advantages:
 - m Simple, inexpensive device
 - m Multi-tier provides graceful degradation: portions of the LAN continue to operate if one of the hubs malfunction
 - m Extends maximum distance between node pairs (100m per Hub)

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Hubs (more)

- r Hub Limitations:
 - m Single collision domain results in no increase in max throughput; the multi-tier throughput same as the the single segment throughput

 - m Individual LAN restrictions pose limits on the number of nodes in the same collision domain (thus, per Hub); and on the total allowed geographical coverage

 - m May not connect different Ethernet types (e.g., 10BaseT and 100baseT)

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Bridges

- r **Link Layer devices:** they operate on Ethernet frames, examining the frame header and selectively forwarding a frame based on its destination
- r Bridge **isolates collision** domains since it buffers frames
- r When a frame is to be forwarded on a segment, the bridge uses CSMA/CD to access the segment and transmit

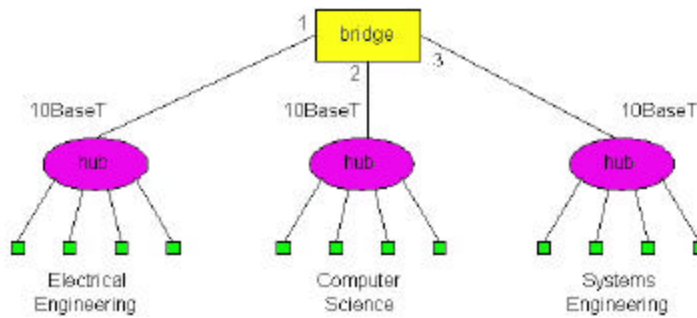
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Bridges (more)

- r Bridge advantages:
 - m Isolates collision domains resulting in higher total max throughput, and does not limit the number of nodes nor geographical coverage
 - m Can connect different type Ethernet since it is a store and forward device
 - m Transparent: no need for any change to hosts LAN adapters

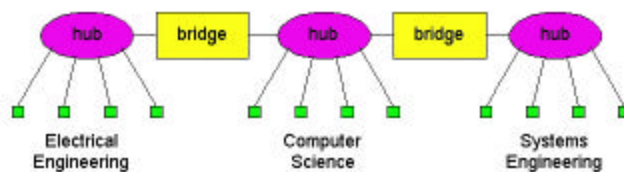
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Backbone Bridge



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Interconnection Without Backbone



- r **Not recommended** for two reasons:
 - Single point of failure at Computer Science hub
 - All traffic between EE and SE must path over CS segment

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Bridge Filtering

- r Bridges learn which hosts can be reached through which interfaces and maintain filtering tables
- r A filtering table entry:
(Node LAN Address, Bridge Interface, Time Stamp)
- r Filtering procedure:
 - if destination is on LAN on which frame was received
 - then drop the frame
 - else { lookup filtering table
 - if entry found for destination
 - then forward the frame on interface indicated;
 - else flood; */* forward on all but the interface on which the frame arrived*/*

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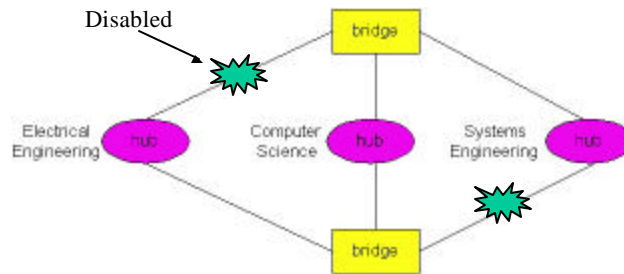
Bridge Learning

- r When a frame is received, the bridge "learns" from the source address and updates its filtering table (Node LAN Address, Bridge Interface, Time Stamp)
- r Stale entries in the Filtering Table are dropped (TTL can be 60 minutes)

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Bridges Spanning Tree

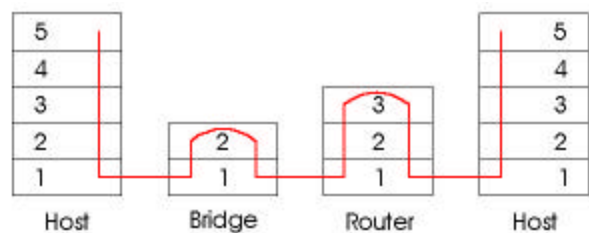
- r For increased reliability, it is desirable to have redundant, alternate paths from a source to a destination
- r With multiple simultaneous paths however, cycles result on which bridges may multiply and forward a frame forever
- r Solution is organizing the set of bridges in a spanning tree by disabling a subset of the interfaces in the bridges:



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WWF Bridges vs. Routers Smackdown

- r Both are store-and-forward devices, but Routers are Network Layer devices (examine network layer headers) and Bridges are Link Layer devices
- r Routers maintain routing tables and implement routing algorithms, bridges maintain filtering tables and implement filtering, learning and spanning tree algorithms



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Routers vs. Bridges

- r Bridges + and -
- + Bridge operation is simpler requiring less processing bandwidth
- Topologies are restricted with bridges: a spanning tree must be built to avoid cycles
- Bridges do not offer protection from broadcast storms (endless broadcasting by a host will be forwarded by a bridge)

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Routers vs. Bridges

- r Routers + and -
- + Arbitrary topologies can be supported, cycling is limited by TTL counters (and good routing prots)
- + Provide firewall protection against broadcast storms
- Require IP address configuration (not plug and play)
- Require higher processing bandwidth
- r Bridges do well in small (few hundred hosts) while routers are required in large networks (thousands of hosts)

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Ethernet Switches

- r A switch is a device that incorporates bridge functions as well as point-to-point 'dedicated connections'
- r A host attached to a switch via a dedicated point-to-point connection; will always sense the medium as idle; no collisions ever!
- r Ethernet Switches provide a combinations of shared/dedicated, 10/100/1000 Mbps connections

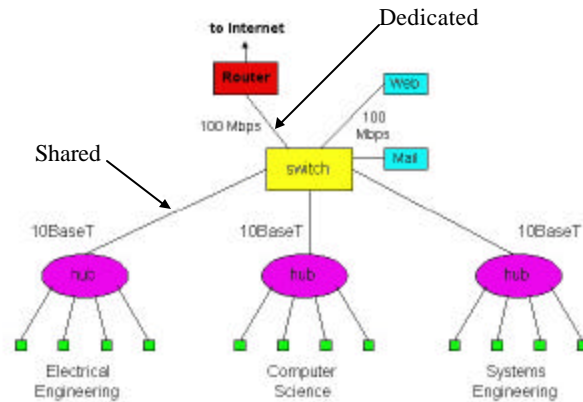
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Ethernet

- r Some E-net switches support cut-through switching: frame forwarded immediately to destination without awaiting for assembly of the entire frame in the switch buffer; slight reduction in latency
- r Ethernet switches vary in size, with the largest ones incorporating a high bandwidth interconnection network

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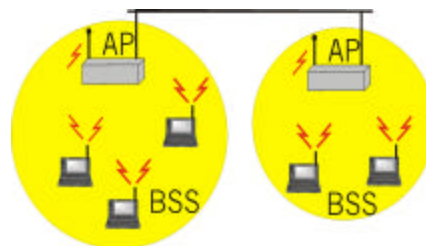
Ethernet Switches (more)



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IEEE 802.11 Wireless LAN

- r Wireless LANs are becoming popular for mobile Internet access
- r Applications: nomadic Internet access, portable computing, ad hoc networking (multihopping)
- r IEEE 802.11 standards defines MAC protocol; unlicensed frequency spectrum bands: 900Mhz, 2.4Ghz
- r **Basic Service Sets + Access Points => Distribution System**
- r Like a bridged LAN (flat MAC address)



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Ad Hoc Networks

- r IEEE 802.11 stations can dynamically form a group without AP
- r Ad Hoc Network: no pre-existing infrastructure
- r Applications: "laptop" meeting in conference room, car, airport; interconnection of "personal" devices (see bluetooth.com); battelfield; pervasive computing (smart spaces)
- r IETF MANET (Mobile Ad hoc Networks) working group



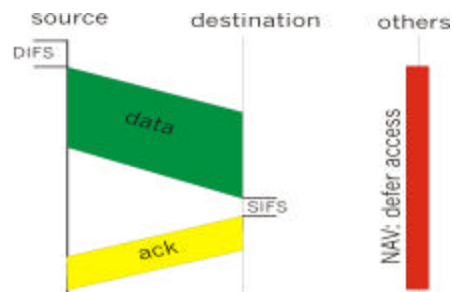
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IEEE 802.11 MAC Protocol

CSMA Protocol:

- sense channel idle for **DIFS** sec (Distributed Inter Frame Space)
 - transmit frame (no Collision Detection)
 - receiver returns **ACK** after **SIFS** (Short Inter Frame Space)
- if channel sensed busy then binary backoff

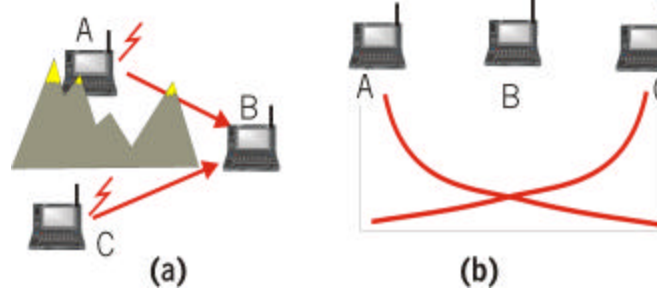
NAV: Network Allocation Vector
(min time of deferral)



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Hidden Terminal effect

- r CSMA inefficient in presence of hidden terminals
- r Hidden terminals: A and B cannot hear each other because of obstacles or signal attenuation; so, their packets collide at B
- r Solution? **CSMA/CA**
- r **CA** = Collision Avoidance

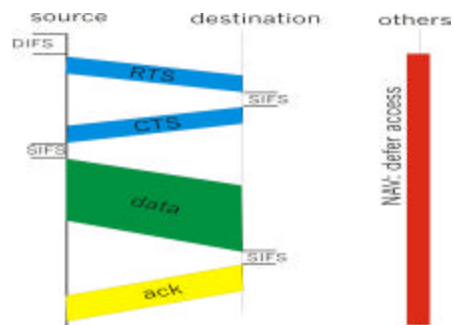


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Collision Avoidance: RTS-CTS exchange

- **CTS** "freezes" stations within range of receiver (but possibly hidden from transmitter); this prevents collisions by hidden station during data
- **RTS** and **CTS** are very short: collisions during data phase are thus very unlikely (the end result is similar to Collision Detection)

•Note: IEEE 802.11 allows **CSMA**, **CSMA/CA** and "polling" from AP



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Point to Point protocol (PPP)

- r Point to point, wired data link easier to manage than broadcast link: no Media Access Control
- r Several Data Link Protocols: PPP, HDLC, SDLC, Alternating Bit protocol, etc
- r PPP (Point to Point Protocol) is very popular: used in dial up connection between residential Host and ISP; on SONET/SDH connections, etc
- r PPP is extremely simple (the simplest in the Data Link protocol family) and very streamlined

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PPP Requirements

- r Pkt framing: encapsulation of packets
- r bit transparency: must carry any bit pattern in the data field
- r error detection (no correction)
- r multiple network layer protocols
- r connection liveness
- r Network Layer Address negotiation: Hosts/nodes across the link must learn/configure each other's network address

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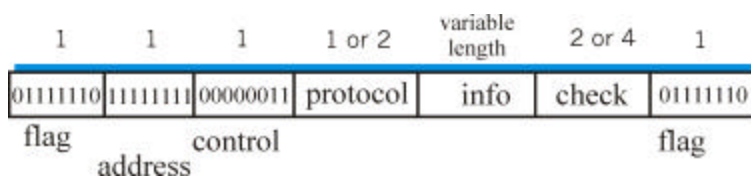
Not Provided by PPP

- r error correction/recovery
- r flow control
- r sequencing
- r multipoint links (e.g., polling)

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PPP Data Frame

- r Flag: delimiter (framing)
- r Address: does nothing (only one option)
- r Control: does nothing; in the future possible multiple control fields
- r Protocol: upper layer to which frame must be delivered (eg, PPP-LCP, IP, IPCP, etc)



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Byte Stuffing

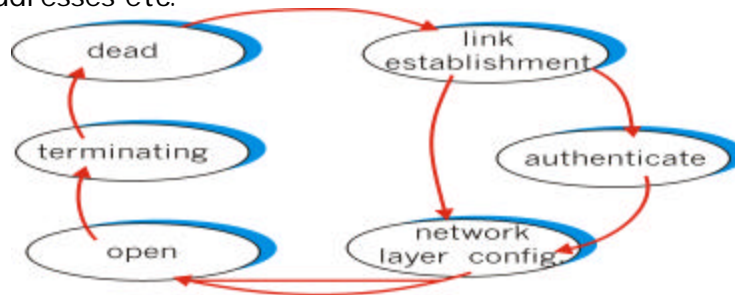
- r For "data transparency", the data field must be allowed to include the pattern <01111110> ; ie, this must not be interpreted as a flag
- r to alert the receiver, the transmitter "stuffs" an extra < 01111110> byte after each < 01111110> data byte
- r the receiver discards each 01111110 followed by another 01111110, and continues data reception



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PPP Data Control Protocol

- r PPP-LCP establishes/releases the PPP connection; negotiates options
- r Starts in DEAD state
- r Options: max frame length; authentication protocol
- r Once PPP link established, IPCP (Control Protocol) moves in (on top of PPP) to configure IP network addresses etc.



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